

KEY-AND-ARGUMENT- UPDATABLE QA-NIZKS

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REMINDER: NIZK

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REMINDER: SIMULATION

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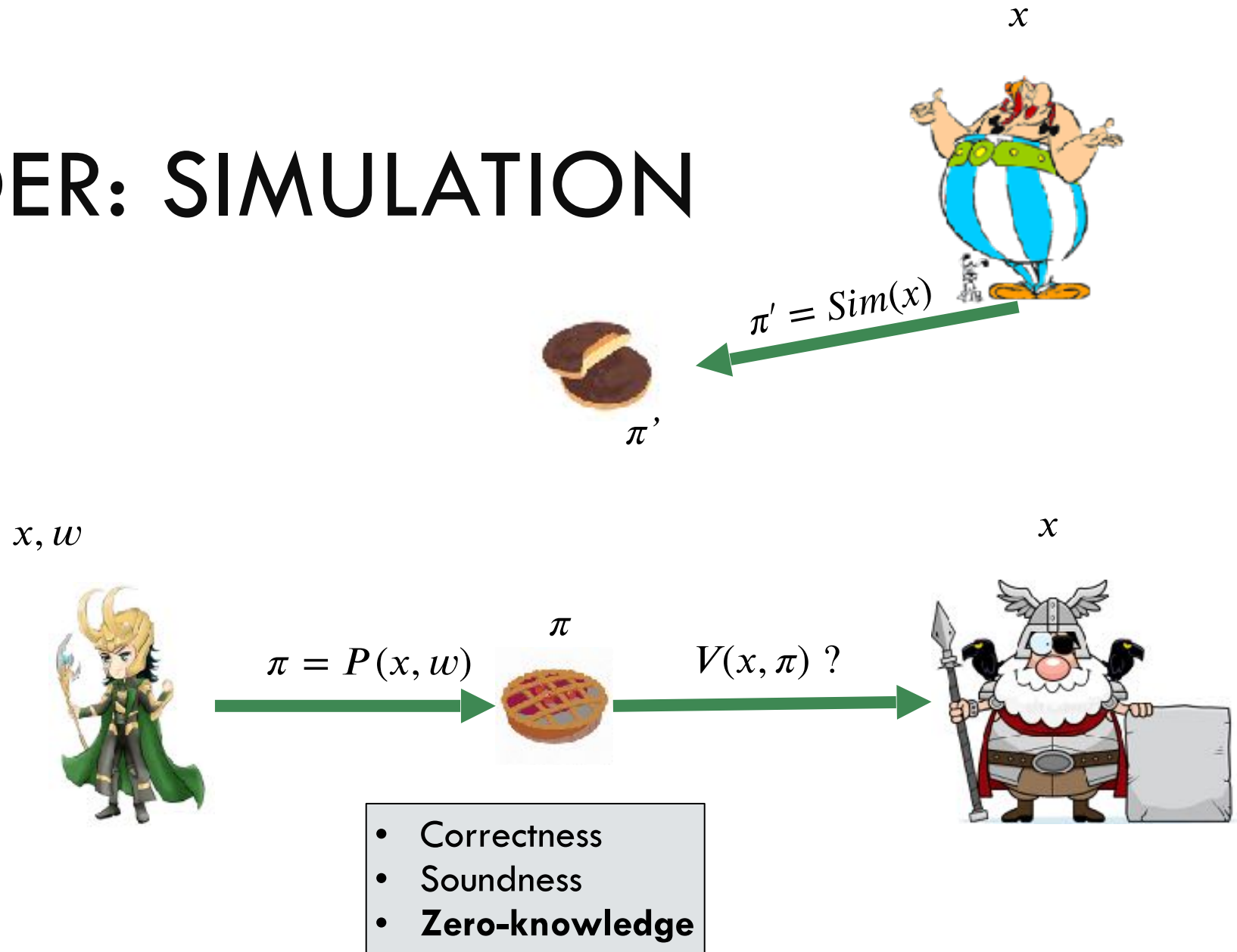
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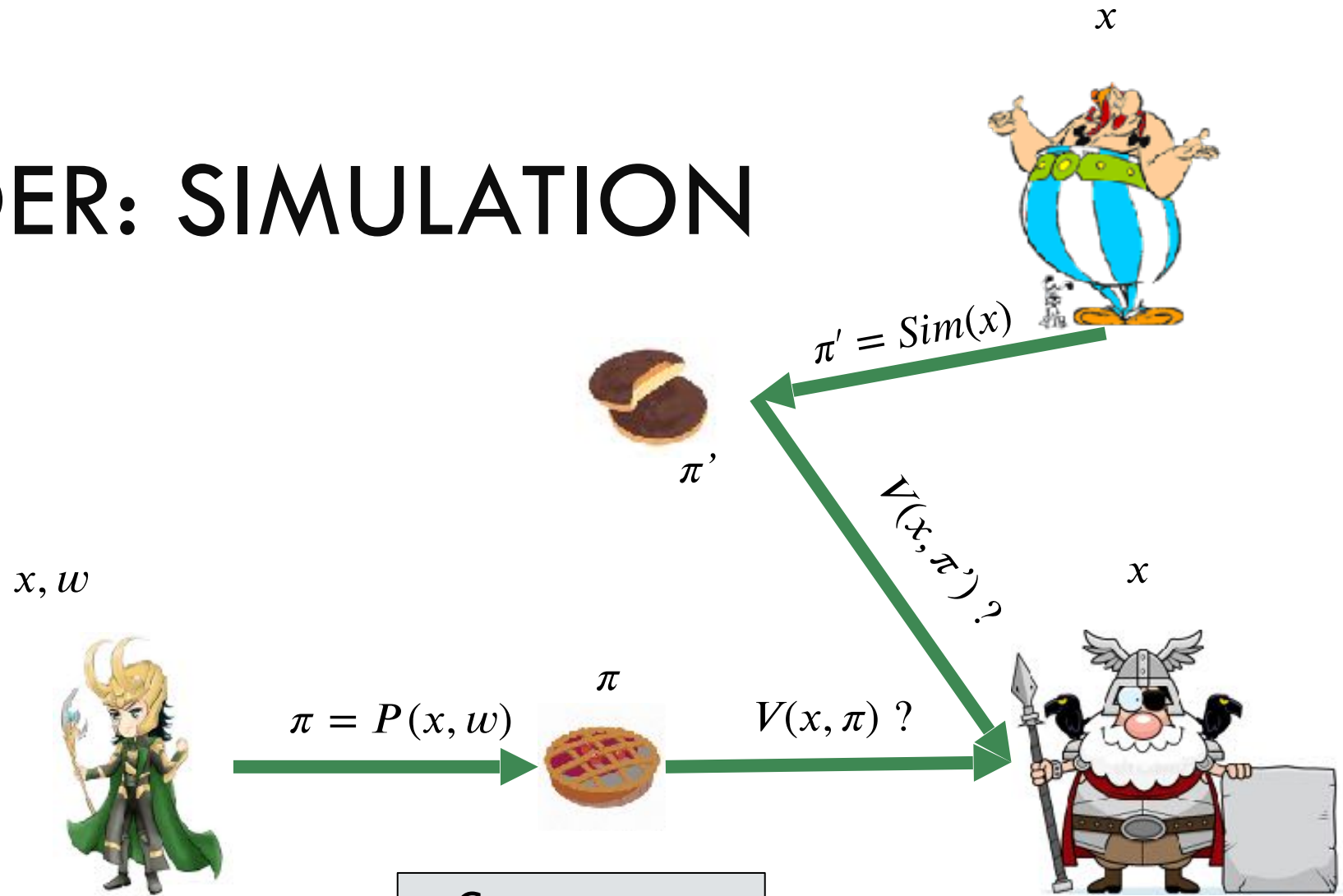


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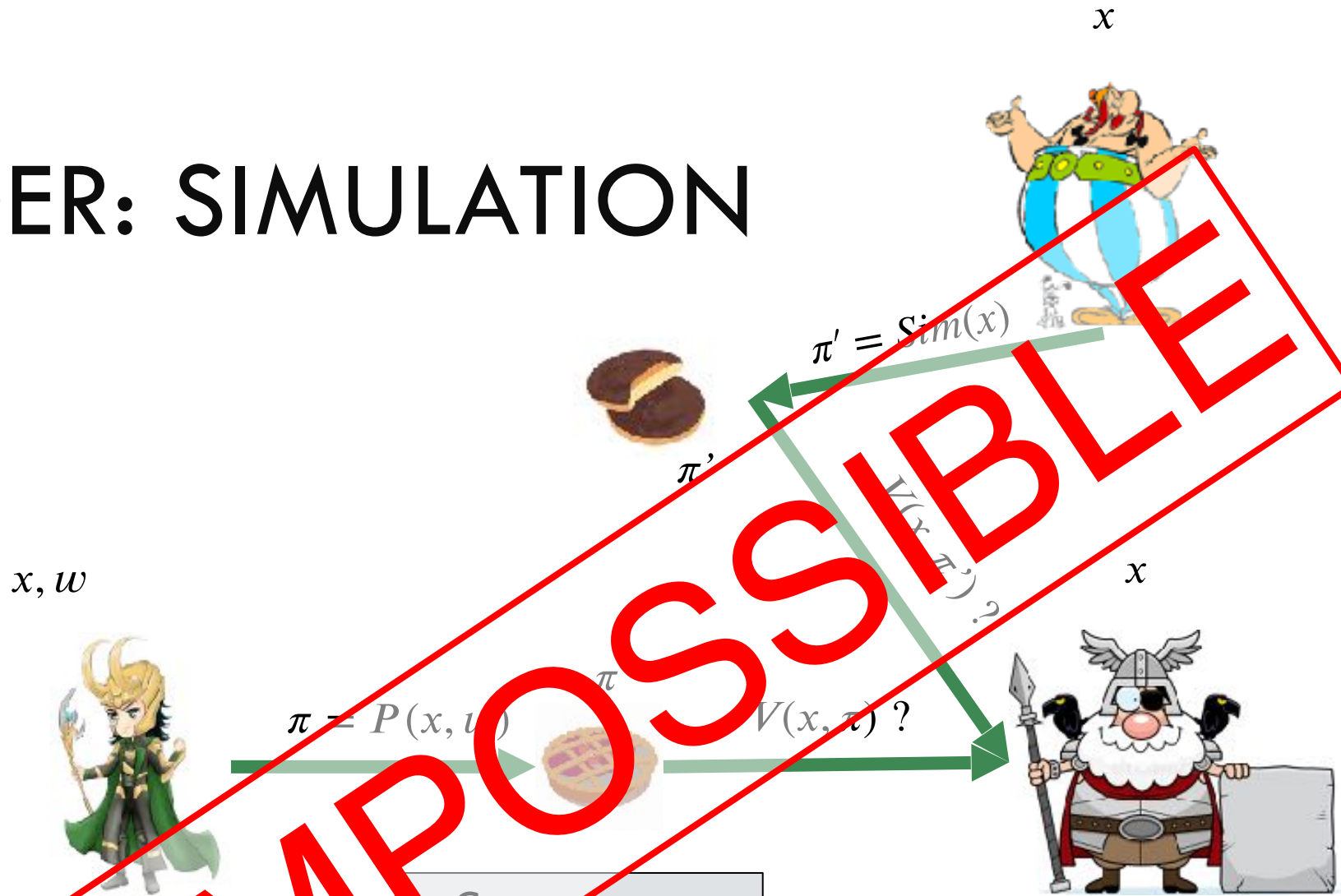


REMINDER: SIMULATION



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IMPOSSIBLE

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crs, x, w



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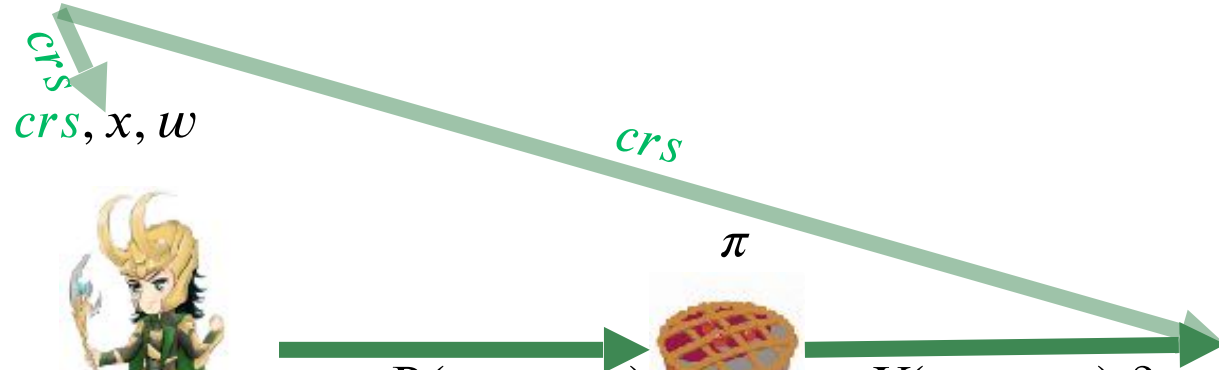


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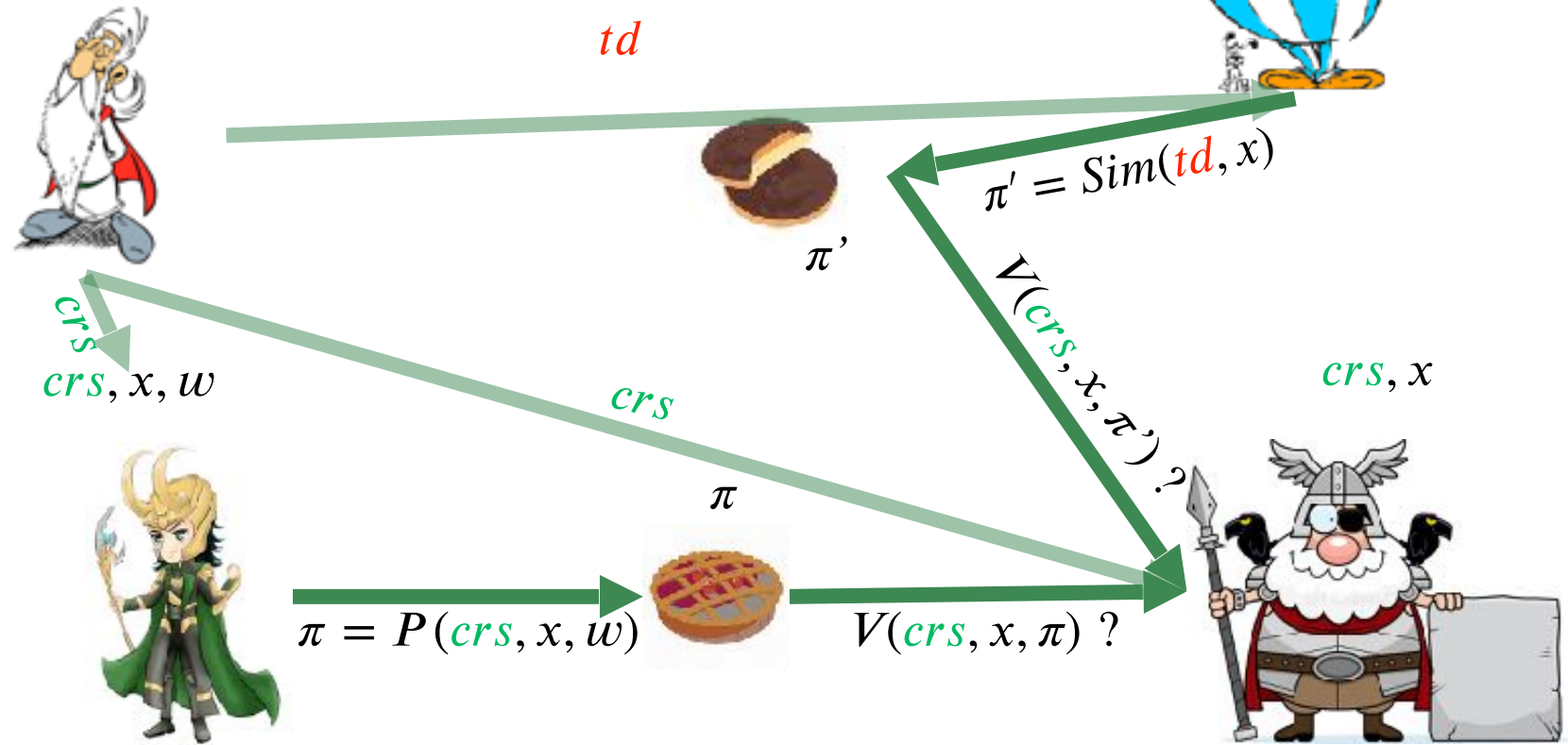
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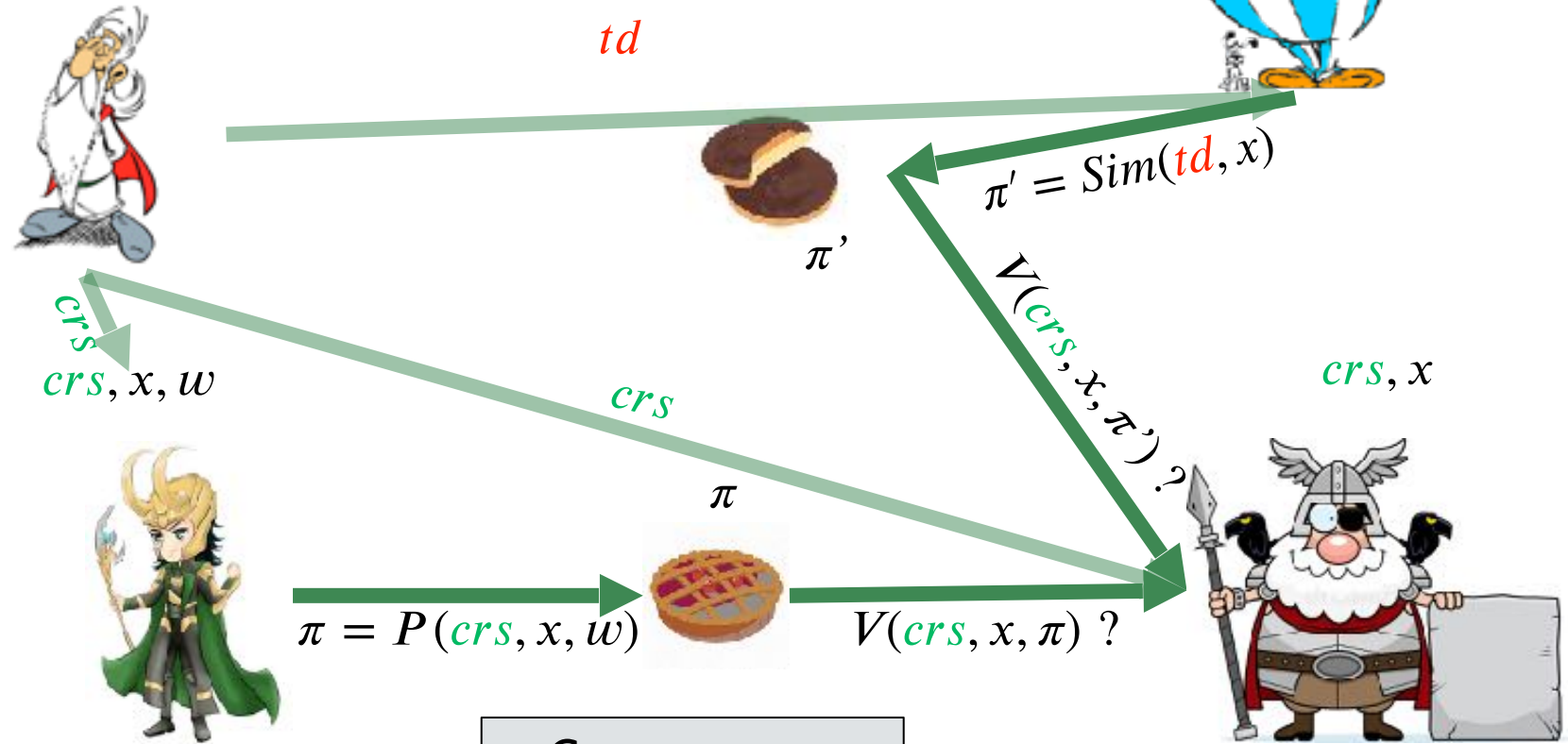
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 - Trust – who is the CRS creator?

EFFICIENT NIZKS (1): ZK-SNARKS

Model	Paper	Trust in CRS creator	Efficiency	Assumptions
CRS Model	Groth10	Everybody trusts	Efficient for NP [GGPR13, Gro16, ...]	Non-black-box for soundness (needed)

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Updatable	GKM+18	P does not trust, V trusts one updater	-..-	-..-

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- building block of SNARK variations

QA-NIZK: DEFINITION

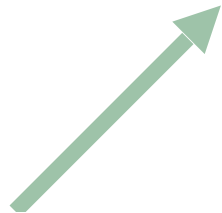
lpar



Usually trusted

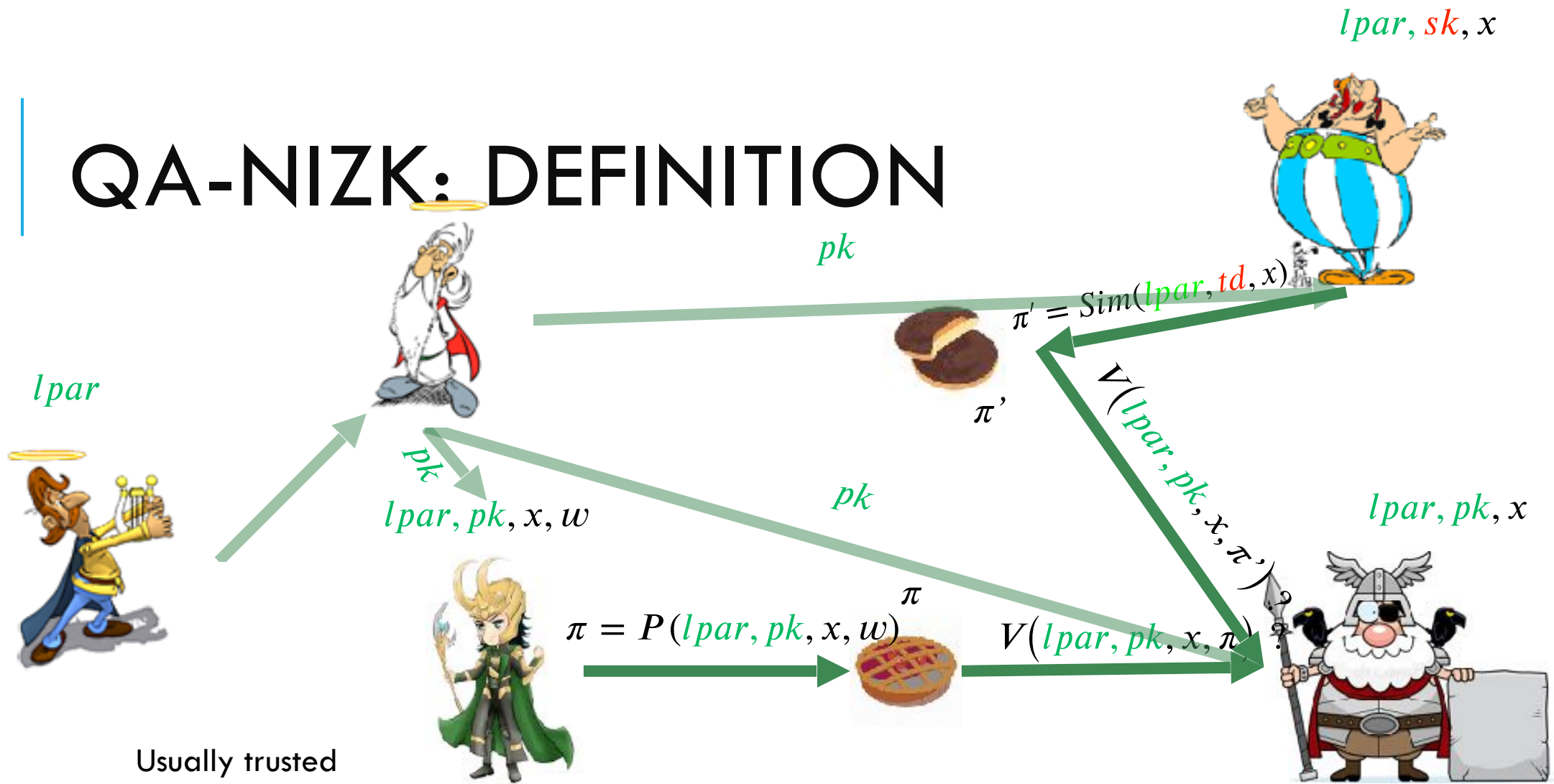
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quasi-adaptiveness enables to get improved efficiency
 Subspace QA-NIZKs where $|\pi| = 1$ group element

EFFICIENT NIZKS (2): QUASI-ADAPTIVE NIZKS

Model	Year	Trust in Ipar &/ pk	Efficiency	Assumptions
CRS Model	JR13	Everybody trusts both	Efficient for linear languages, less efficient for NP	Black-box
Sub-ZK (BPK model)	ALSZ20	V trusts pk generator	-..-	NBB for ZK (needed)
Updatable	This work	V trusts <u>one</u> pk updater	-..-	-..-

UPDATABLE SNARKS: IDEA

- **Non-updatable zk-SNARKs**

- $sk = \mathbf{x}$ — a vector of uniformly random integers

- $pk_C(\mathbf{x}) = \{[f(\mathbf{x})] : f \in S_C\}$ for a **C-dependent** function set S_C

- **Updatable SNARKs**

- $pk_u(\mathbf{x}) := \{[f(\mathbf{x})] : f \in S_n\}$ for a universal set S_n

- ...universal for all circuits of size $\leq n$

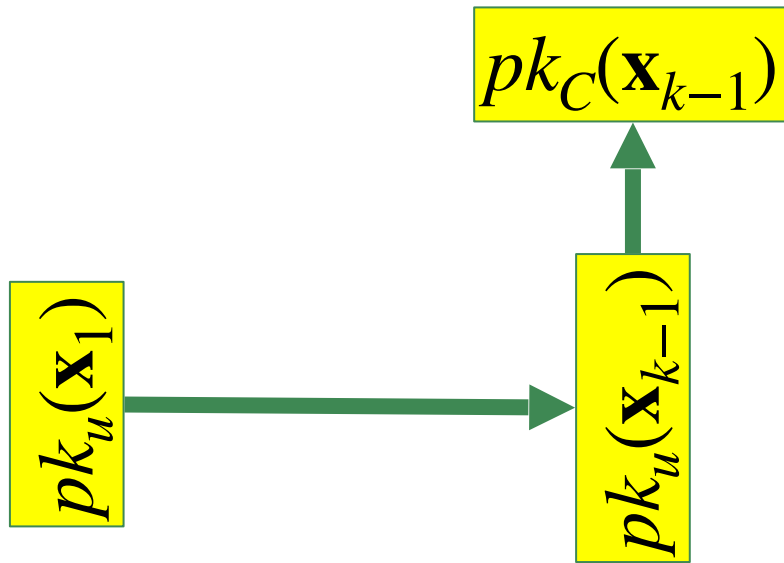
- verifiably specialize $pk_u(\mathbf{x}) \Rightarrow pk_C(\mathbf{x})$ to concrete circuit C

UPDATABLE SNARKS: DETAILS

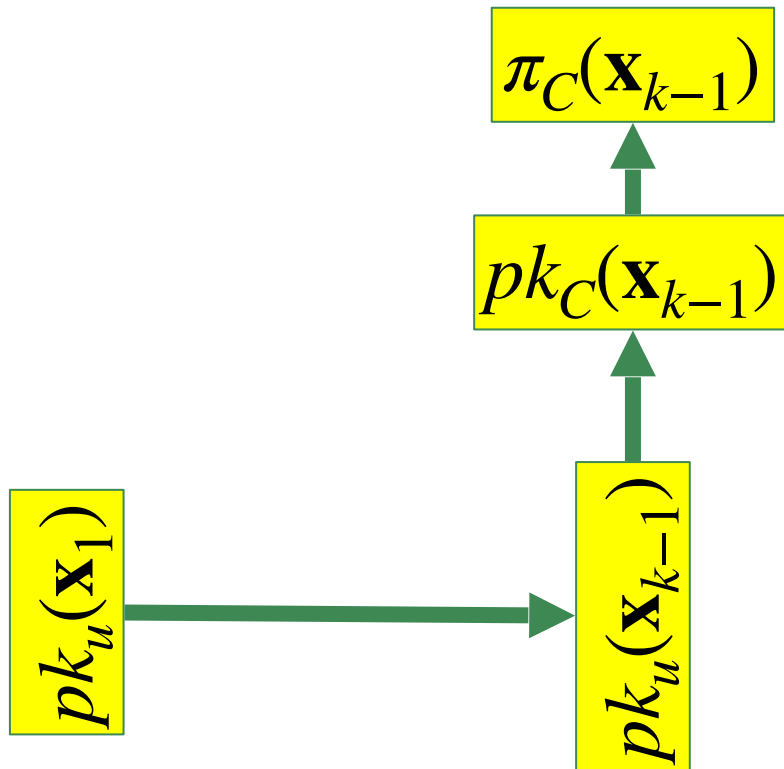
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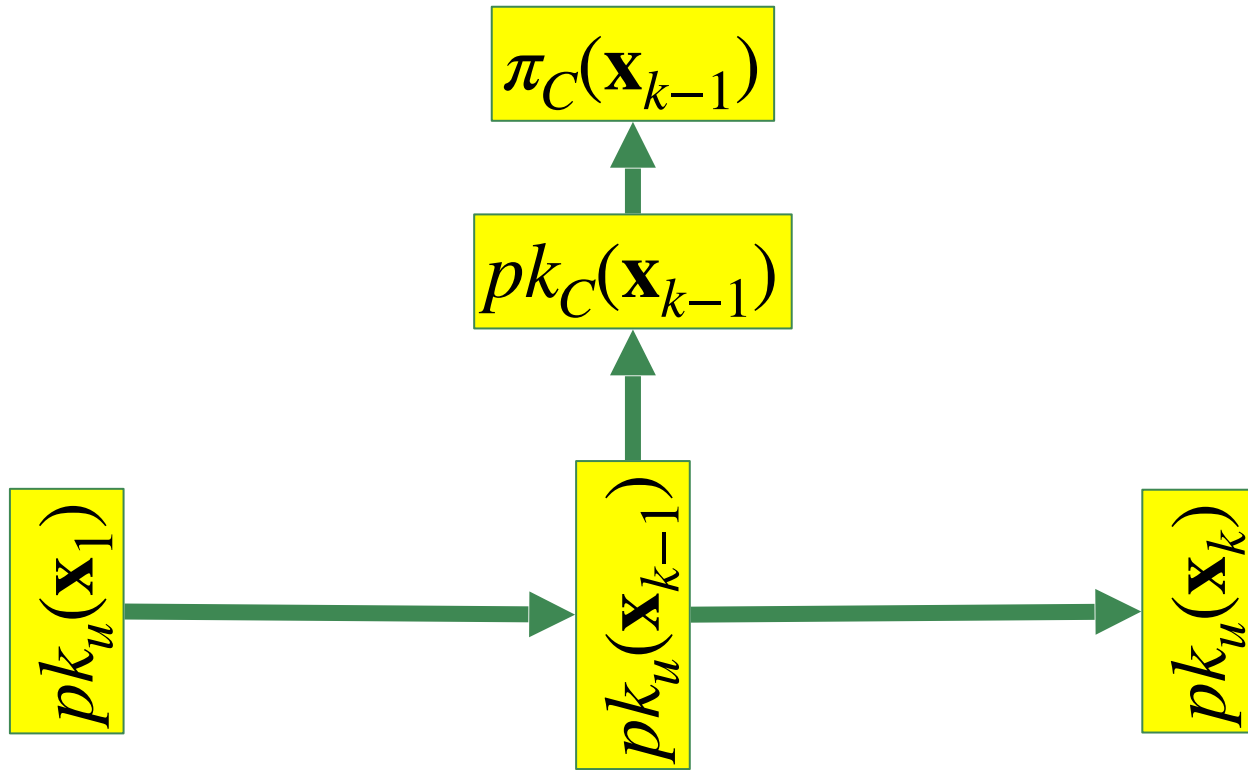
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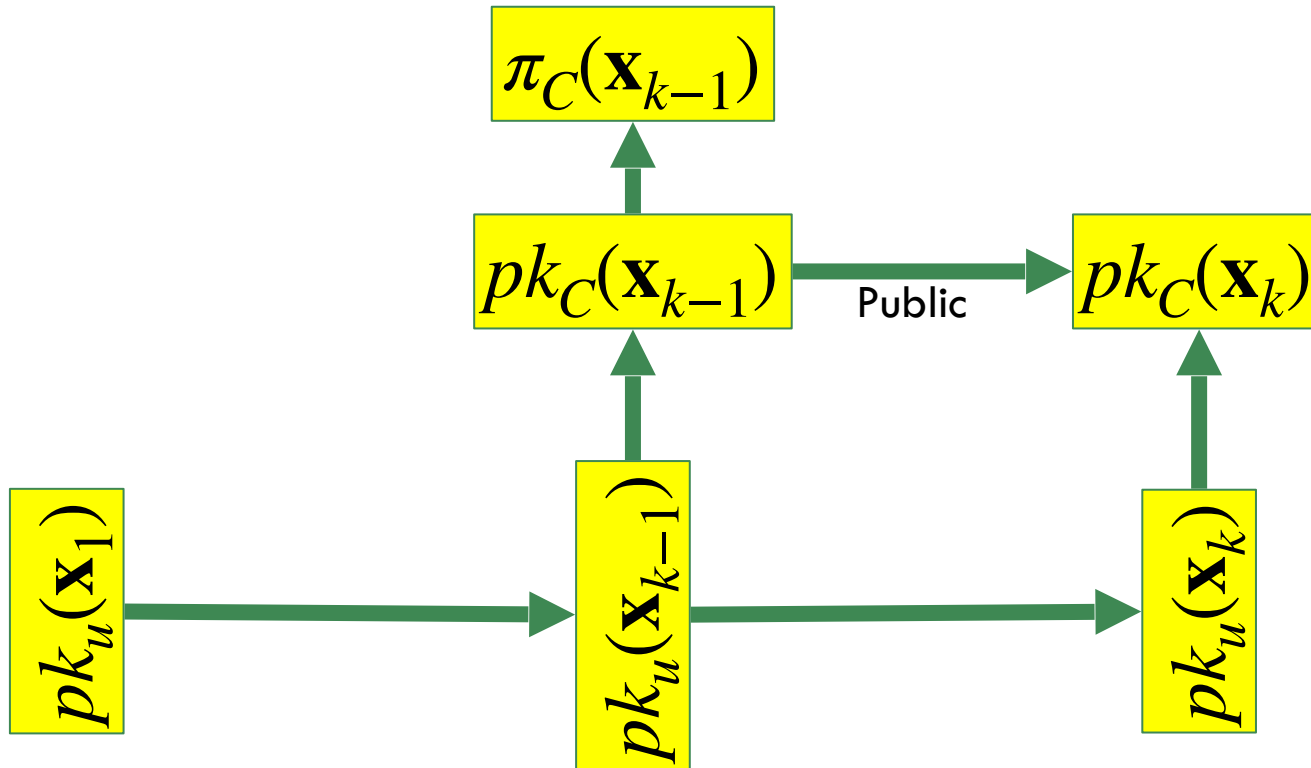
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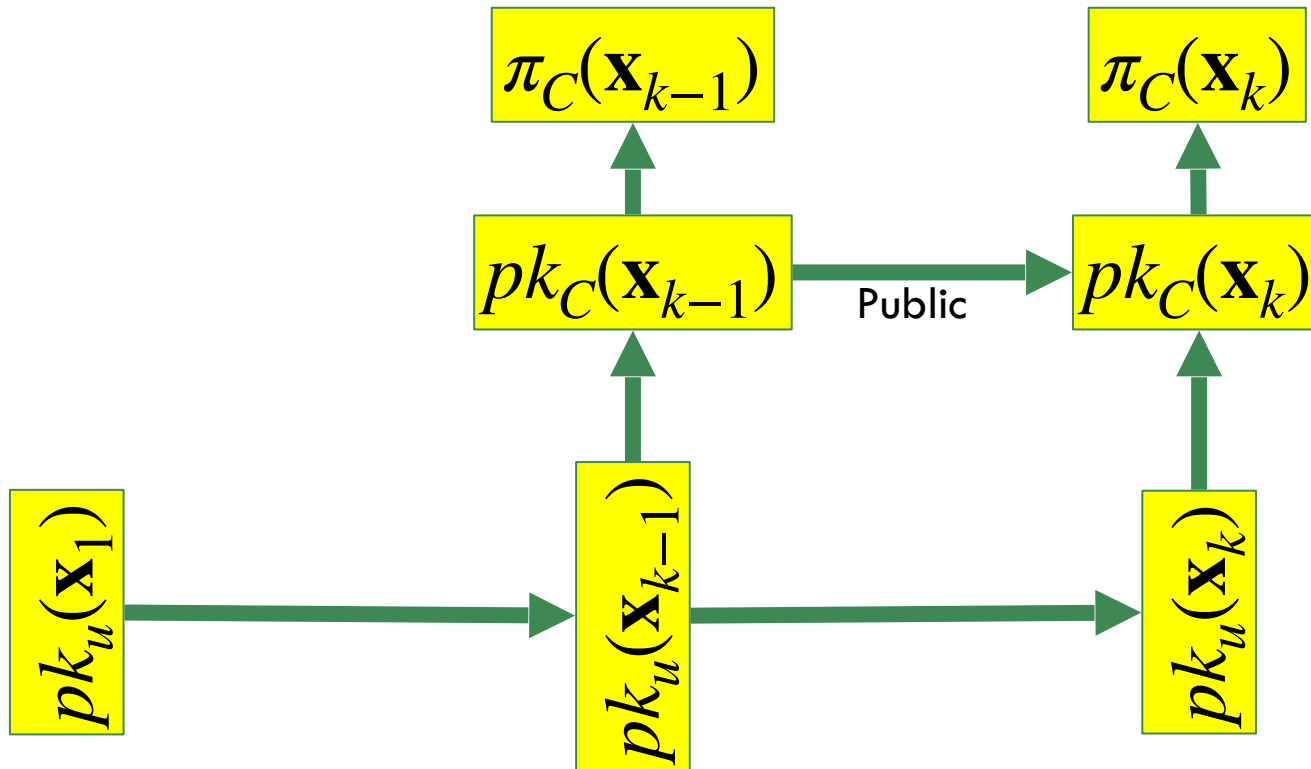
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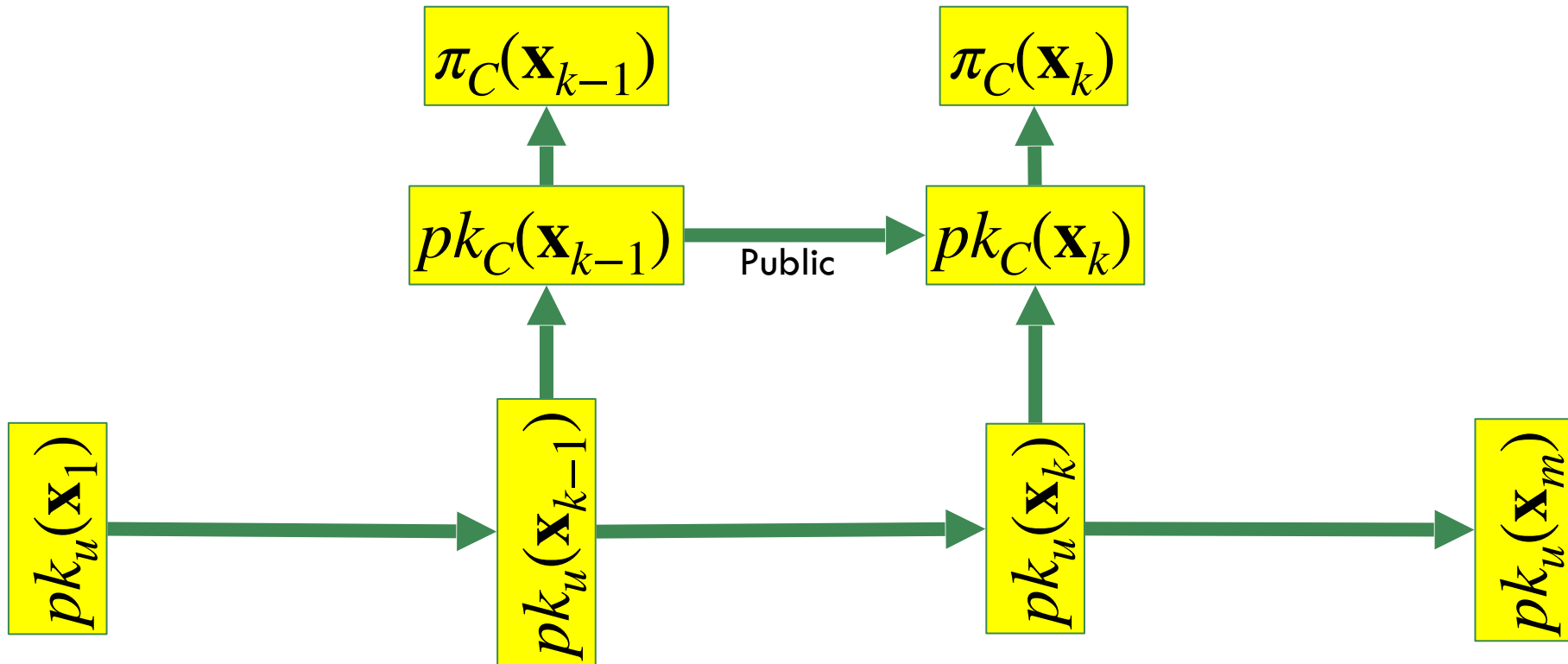
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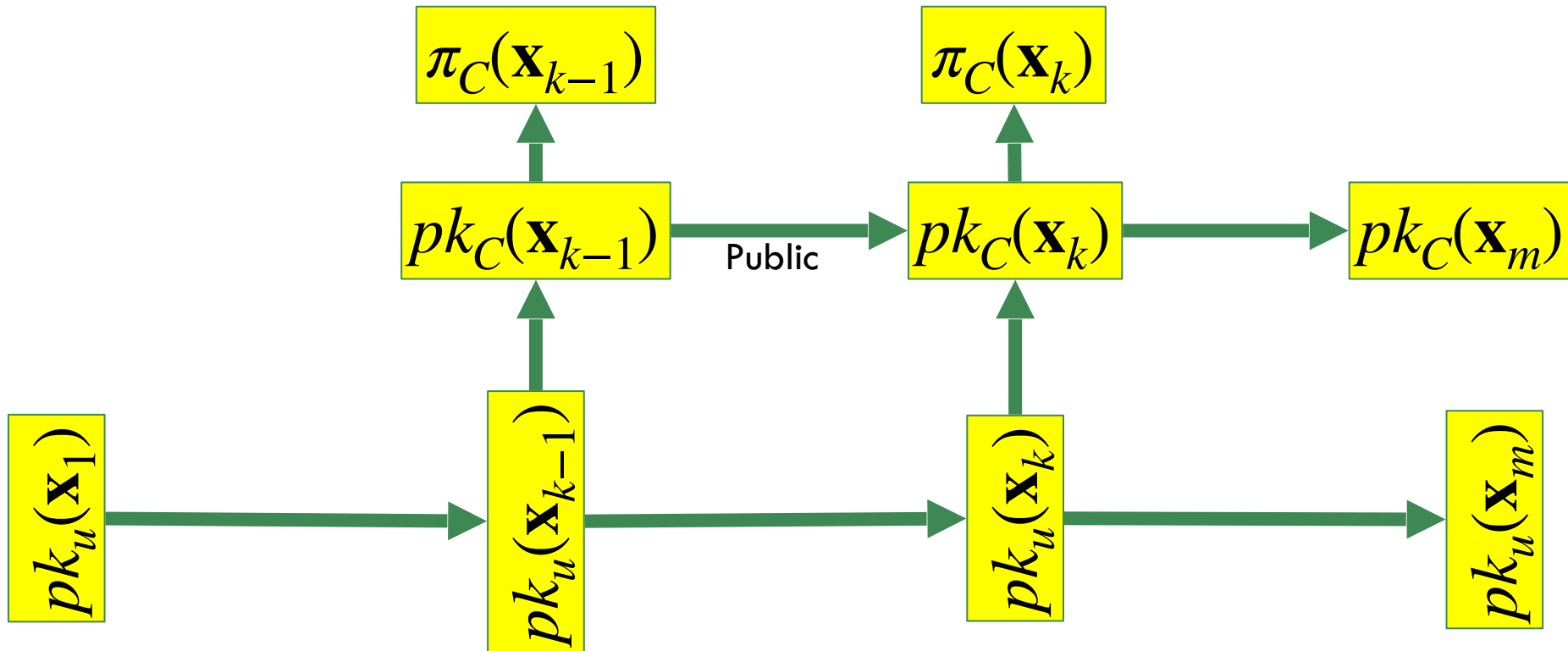
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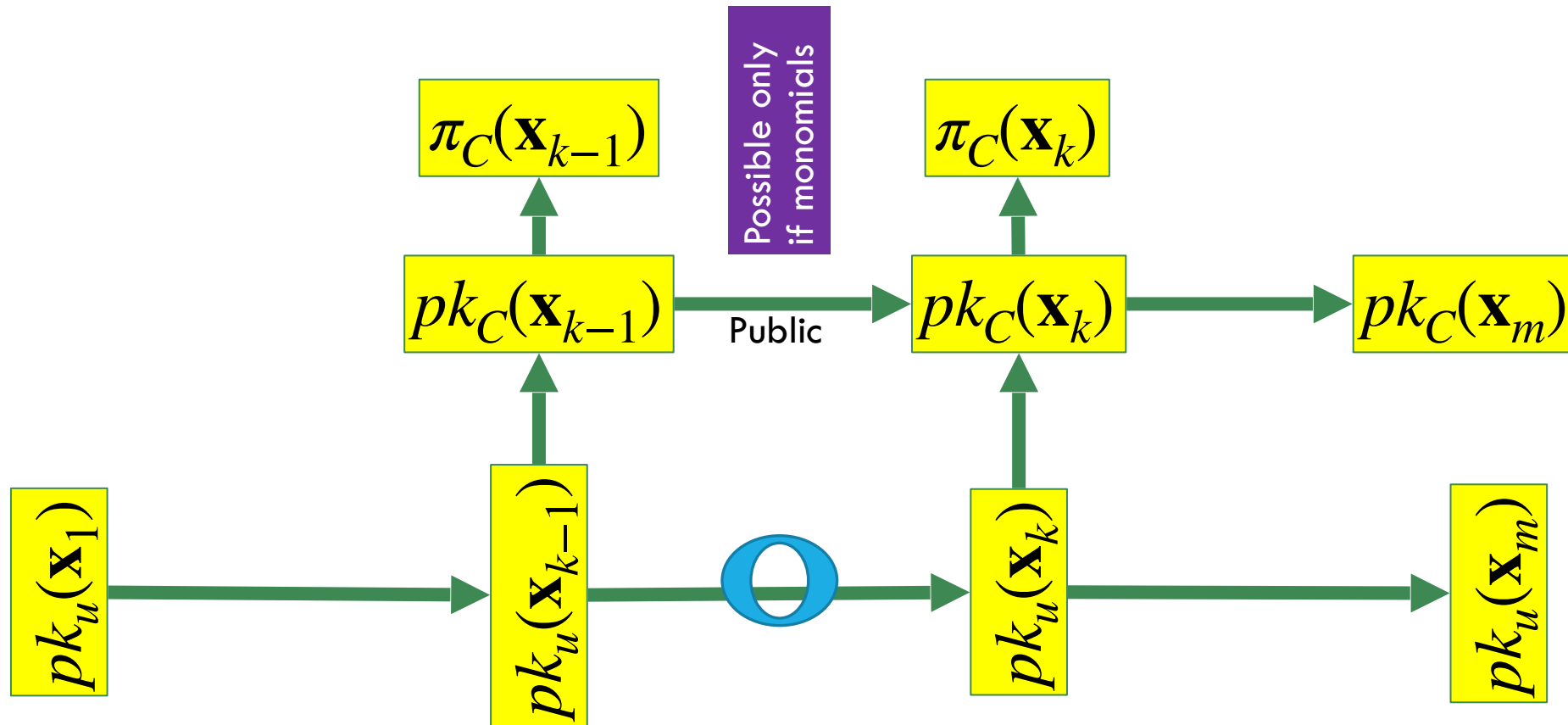
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 - Our observation: for this we need the distribution of sk belongs to a generalized ideal of a convolution semigroup (“nice distribution”)

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 - ... cannot reuse existing techniques of [GGPR13, Gro16, ...]

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 - Update the rest of the pk separately in each application

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Caveat: updated π is as sound as original π

